DISTRIBUTED AGENT-BASED ONLINE AUCTION SYSTEM

Costin Bădică, Sorin Ilie, Alex Muscar
University of Craiova, Software Engineering Department
Bvd. Decebal, 107, Craiova, Romania
e-mail: {cbadica, silie, amuscar}@software.ucv.ro

Amelia Bădică
University of Craiova, Statistics and Business Information Systems Department
Str. A. I. Cuza, 13, Craiova, Romania
e-mail: ameliabd@yahoo.com

Liviu Sandu, Raluca Sbora
University of Craiova, Software Engineering Department, Bvd. Decebal, 107 Craiova, Romania
e-mail: {ssanduliviu, ralus08}@yahoo.com

Maria Ganzha
Polish Academy of Sciences, Systems Research Institute, ul. Newelska 6
01-447 Warszawa, Poland
&
University of Gdańsk, Institute of Informatics, ul. Wita Stwosza 57
Gdańsk, Poland
e-mail: Maria.Ganzha@ibspan.waw.pl

Marcin Paprzycki
Polish Academy of Sciences, Systems Research Institute, ul. Newelska 6
01-447 Warszawa, Poland
&
Warsaw Management Academy, ul. Kaweczynska 36
Warsaw, Poland
e-mail: Marcin.Paprzycki@ibspan.waw.pl
Abstract. This paper concerns the design and development of a distributed agent-based online system for English auctions. The proposed system is composed of two parts: an Agent-based Auction Server and a Web-based Graphical User Interface. The first part of our work brought about the advantages introduced by the multi-agent systems technology to the high-level of abstraction, modularity and performance of the server architecture and its implementation. On the server side, bids submitted by auction participants are handled by a hierarchical organization of agents that can be efficiently distributed on a computer network. This approach avoids the bottlenecks of bid processing that might occur during periods of heavy bidding, like for example snipping. We present experimental results that show a significant improvement of the server throughput compared with the architecture where a single auction manager agent is used for coordinating the participants for each active auction that is registered with the server. The second part of our work involved analysis of external functionalities, implementation and usability of a prototype online auction system that incorporates the Agent-based Auction Server. Our solution is outlined in terms of information flow management and its relation to the functionalities of the system. The main outcome of this part of the work is a clean specification of the information exchanges between the agent and non-agent software components of the system. Special attention is also given to the interoperability, understood here as successful integration of the different data communication protocols and software technologies that we employed for the implementation of the system.

Keywords: Distributed system, multi-agent middleware, online auction

1 INTRODUCTION

The vision of e-commerce automation proposes the development of global e-commerce environments populated by software agents, thus enabling the dynamic development of trading relationships between business partners. In particular, increasing the level of automation of negotiations is needed, to allow the engagement of stakeholders, either individuals or business organizations, into nontrivial dynamic business relationships.

As argued in [8, 16], auctions provide a general solution to the problem of discrete resource allocation among selfish agents in a multi-agent system. There are many types of auctions including single-good, multi-unit, combinatorial auctions, and double auctions. Auctions represent a special class of negotiations with many applications in conducting e-business transactions [17, 36]. In particular, auctions are useful for trading in the following areas: spectrum licenses, electricity markets, emission rights, airports takeoff and landing slots, exploitation rights of natural resources (e.g. oil-drilling), selling of collectibles, antiques, luxury and second-hand products, government procurement contracts, foreign exchange, a.o. [7].
As online auctions spread with the advent of the web, many types of online applications for auctions were proposed, including auction directories, auction tops, meta-auctions, and auction servers [2]. Recently, the research focus was set on the development of more process-generic, flexible and reusable auction solutions, with an increased potential for applicability both to the B2C and B2B sectors. In this context, application of agent-based systems was proposed as a new approach that takes the idea of an auction service from the human-driven web to the software agents’ world [13, 18].

In this paper we consider specific example of English auctions. In this case, the seller (or the auctioneer that represents the seller) announces an initial price for the goods (assuming an e-commerce application setting) and auction participants bid increasing amounts of money during a predefined time frame, usually by some minimum increment set by the seller. At the end of the bidding process, the agent with the highest bid is declared the winner of the English auction. Note that we consider here the deadline-driven model of the English auction, rather than the time of inactivity model.

Our current work in this area is focused on the analysis, design and implementation of an open, flexible infrastructure for agent-based automated negotiations. The two objectives of this paper are:

1. To underline the advantages of using multi-agent systems and state-of-the-art agent frameworks and middleware [29, 19] when developing a realistic auction server. Here, the focus is on employing clean software engineering principles (abstraction and modularity), as well as on evaluating and improving the performance and scalability of the implementation. Initial discussion of software engineering principles that underline this work, as well as preliminary performance assessment were presented in [2]. The scalability and performance aspects were further expanded by realization of a cluster-based implementation that was outlined and initially evaluated in [11]. Here, we combine and further extend the results presented in these two conference papers.

2. To develop a tool that can be used for online auctions in B2C systems. This goal is addressed by focusing on the details of incorporating the Agent-based Auction Server into a Web-based application for online auctions.

The proposed agent-based solution for the auction server combines the best features of:

1. generic software framework for automated negotiations [3];
2. market architecture for auction development [6];
3. rule-based declarative representation of auction mechanisms [1, 4];
4. special computing nodes available in active networks and realized by means of proxy agents [9];
5. agent-based service-oriented architecture [2].
Consequently, it provides certain features including openness, generality, and scalability. For example, with this approach, an auction is seen as a separate service, rather than being entirely incorporated into an e-shop infrastructure. This can be seen as a gain in openness, as the service is now open for rental and configuration by the e-shop that would like to sell its products through an auction, for example for clearing its shelves during the “sales” time. Generality comes from the fact that the e-shop can now choose the most appropriate auction server depending on factors like performance, reliability or trustworthiness of the service. Finally, scalability comes from our new approach that combines the use of Proxy agents with two level balanced tree structures for handling participants’ bids.

The second part of our work was focused on the analysis of the external functionalities, implementation and usability of a prototype online auction system that incorporates the Agent-based Auction Server. We present the details of our solution in terms of the information flow management and its relation to the functionalities of a system for online auctions.

In summary, our work brings several contributions to the research on agent-based e-commerce.

When an auction (in particular an English auction) is modeled as a multi-agent system, the agents’ interaction through message exchanges is required for the bidding and price update activities. In our previous approach [2], the central agent for auction management was responsible with handling all the communication with, and between, the auction participants. So, if a large number of participants joined the auction, this agent became highly stressed. Second, although the approach introduced in [2] was described as distributed, understood as possibility for agents to be arbitrarily located on networked computers, it was still biased towards centralization, because the central agent was a system bottleneck, and thus hindered the scalability.

In this paper we propose a solution aiming at improving the architecture of our Agent-based Auction Server by relaxing the central agent from part of the stress caused by its heavy load of message handling. Based on the idea that was initially proposed by [9], for improving the performance of on-line auction systems using special computing nodes available in active networks, we enhanced our system with the introduction of third party agents called proxies. Each Proxy agent will handle a part of the communication with the auction participants.

So, while in our prior approach the Auction Manager would receive all the bids from the Participants, with this new approach the Participants are split into disjoint groups, and each group is managed by a single Proxy agent. The Participants communicate heavily with their proxies, while proxies pass on to the Auction Manager agent only the relevant bids, while the other bids are filtered out and processed locally, thus reducing the amount of messages handled by the Auction Manager.

Note that the communication between two agents is faster when the agents are located on the same machine, rather than when they are located on separate machines. When our auction server is distributed on several computers, some agents will have to exchange messages over the network, thus increasing the communication.
time, as well as the overall server response time. So, with our solution we also aimed to improve the communication time between agents, when the server is distributed on several computers, by keeping, whenever possible, Proxy agents on the same machine with their “participant agents”.

Finally, our work contributes also to research concerning usability of agent-based e-commerce solutions. While very attractive, the complete automation of e-commerce processes is probably impossible to achieve, and therefore the human user involvement, through an appropriate online system, will be always necessary. Therefore, we found it important to experiment with the integration of the proposed Agent-based Auction Server into a usable online auction system that allows the direct human user involvement in auctions, via a Web-based GUIs.

Furthermore, we provide a clean specification of the information exchanges between the agent and non-agent software components of the system, which is particularly interesting from the software engineering point of view. In this context, special attention is also given to the heterogeneity of the different data communication protocols (for example: HTTP, FIPA, parameter passing via method invocation, a.o.) and software technologies (Web technologies vs software agent technologies) that we utilized for implementing the system.

The paper is structured as follows. We start in Section 2 with a brief overview of related works in the field. In Section 3 we present the architecture of our Agent-based Auction Server covering agent types, interaction protocols and mechanisms for efficient bid processing. Next (in Section 4), we propose the design of a Web system for online auctions that incorporates the auction server. Here, the discussion is focused on three aspects:

1. system architecture,
2. design details of the Web layer as well as of the interfacing of the agent and non-agent software, and
3. interaction protocols.

In Section 5 we present results of experiments carried out with the auction server, including recorded values of latency and throughput parameters. Here we also discuss the usability of the online auction system. In the last Section 6, we present our conclusions and we point to future works.

2 RELATED WORKS

The interest in development of online software systems for online negotiations, with a special focus on online auctions, increased significantly during the last 15 years. Traditionally, auctions were utilized for trading support in economic markets in offline as well as in online environments. Recently auctions started to be applied in market environments for trading resources for utility computing, including grids and clouds [14].
One of the first and most influential works in the area of auction servers for online applications is the Michigan Internet AuctionBot introduced in [26]. This is a versatile and robust server for online auctions supporting both agent-oriented and human-oriented auction execution. The Michigan Internet AuctionBot introduced the principles of software design for supporting flexible auction mechanisms, including: separation of the user interface from the core auction engine, the capability of running multiple auctions concurrently, as well as the abstraction of the auction process. Most of these principles are currently employed by state-of-the-art auction servers including our own.

In [21], the authors proposed an Internet-based negotiation server for e-commerce applications. Although this work does not explicitly address auction mechanisms (rather, the focus is on bargaining) and the use of software agent technologies, it is interesting for our approach for at least the following reasons: i) the system is conceptualized as a replicable service that can be multiply instantiated by complementing standard Web server software, i.e. quite similarly to our proposal; ii) the system incorporates methods of event-based rule processing and constraint satisfaction for checking negotiation proposals and implementation of negotiation strategies which, although they are not the focus of this paper, were also employed in our previous work ([1, 2]).

The authors of [20] propose an agent-based modeling of the New York Stock Exchange specialist system. Although this work clearly differentiates from our own work, as the focus is not on the development of an online system incorporating an auction server, but rather on the agent-based modeling of the complex interactions occurring in the New York Stock Exchange specialist system, there are also similarities. First, their modeling addresses a non-trivial class of auctions – continuous double auctions and, second, the modeling could be further expanded to cover the development of an e-service system as part of the New York Stock Exchange.

The e-Game tool that supports the design and implementation of electronic market simulation games inspired by the real life problems was proposed by the authors of [15]. These simulations can also incorporate various types of auctions, and they were used for teaching purposes. The e-Game tool provides both Web and agent interfaces, similarly to our system. Nevertheless, unlike our work, the aspects related to software engineering principles, performance and scalability were not addressed.

In [23], the authors present the principles of constructing online auction systems that were employed for building the Research Auction Server for performing both simulated and real auctions. However, many details are missing in their description, especially those related to the interaction protocols. Therefore we could not compare our approach with [23] because of the missing information. Moreover, although the “agent” metaphor is used we noticed that the development of the Research Auction Server did not actually use software agent technologies.

A generic online auction server was discussed in [27]. The server supports a flexible bidding language based on the OR/XOR formulae. Although apparently there are many similarities with our own work, the details of the design and implemen-
tation of the system are actually missing; only a listing of available technologies is provided. In particular, the interaction protocols and the details of the interfacing of agent and non-agent software are not described.

A configurable auction server was also proposed by the authors of [24]. This server targets resource allocation in the grid and therefore its design addresses the heterogeneity of the grid environment by allowing the dynamic configuration of the auction mechanism to meet the application requirements.

The authors of [30, 31] propose an agent-based infrastructure for autonomous services for management of the contracting of Cloud resources that covers also negotiation. Their system generates a service-level agreement – SLA representing the result of the resource negotiation and booking with available providers. The use of SLA has the advantage that it can be re-negotiated and monitored – a feature, which is missing in our approach. Moreover, while our work is more suitable for auctions, papers [30, 31] are focused on other negotiation mechanisms, like e.g. Contract-Net [32]. Thus, they are closer to our proposed negotiation framework introduced separately in [33].

3 DISTRIBUTED AGENT-BASED AUCTION SERVER

There are many definitions of the agent concept [34]. For the purpose of this work, by software agent (agent in what follows) we understand a software entity that:

1. has its own thread of control and can decide autonomously if and when to perform a given action;
2. communicates with other agents by asynchronous message passing;
3. can be referenced using its name, also known as agent identifier;
4. can be located on an arbitrary machine in a computer network, providing that a certain runtime environment is locally available.

This runtime is usually known as agent platform (see [29] for a recent overview of agent programming languages and platforms). In our current work, we use the Java Agent DEvelopment Framework – JADE [10] agent platform.

In this section we outline the architecture of our Agent-based Auction Server, highlighting agent types and relationships between them and users. Furthermore, we describe agent interactions:

1. inside the auction server and
2. with external agent and non-agent software.

3.1 Agent Types and Their Functions

Let us now summarize the types of agents included in our auction server, focusing on their functionalities. The initial architecture and the agent interaction protocols
of the server were introduced in [2]. In [11] we proposed an improved architecture that enables the deployment of the server on a computer cluster.

The auction server was designed to support the innovative concept of generic agent-based auction service. It is represented by a collection of cooperating agents that interact inside the server, as well as with its external environment, using agent interaction protocols.

The software infrastructure of the server contains the types of agents depicted in the class diagram in Figure 1. The auction server is actually composed of three main parts or layers: core, resource manager, and interface.

![Figure 1. Relationships between users, agent types and auctions on the auction server](image)

### 3.1.1 Interface Layer

The Personal Agent, Participant, and Initiator Participant agents compose the layer that realizes the interface of the server with its external environment.

The Personal Agent is residing on the server side and it connects the user with the auction server. For each user registered with the server there is exactly one Personal Agent created on the server. This agent gets input from the user, through an external user interface. This can be achieved directly, i.e. the Personal Agent can incorporate a user interface, or via a binding software that connects the Personal Agent with an external Web-based GUI. There is a one-to-one mapping between the user name and the identifier of the corresponding Personal Agent.
Whenever a human user connects to the server, usually using an external program (for example a GUI or a Web browser), a new Personal Agent is eventually created and assigned to the user (if no such agent is already active on the server). We set our server up to be used in a Web environment using the servlet technology on the server side (see Section 4 for details). In this case, it is the responsibility of the servlet to talk to the Personal Agent, when the user logged into the system, via an appropriate binding software.

The Personal Agent allows the user to perform the following operations:

- to create a new auction
- to subscribe to an existing auction
- to submit bids to one of his or her subscribed auctions, via his or her Participant agent
- to receive notifications about status updates of his or her subscribed auctions
- to receive accepts and rejects for his or her submitted bids.

Personal Agents can be importantly enhanced with complex behaviors (like for example behaviors specific to Belief-Desire-Intention (BDI) architecture and agent-programming languages [39]) that would allow their truly autonomous operation for better serving the interests and goals of the human user. However, this path of investigation is outside the scope of this paper and it was left as future work.

Participant agent represents a Personal Agent that serves a user (usually with the buyer or seller role, although an auctioneer role can also be used, for example, in double auctions), registered and engaged into a particular auction. For each user registered to participate in an auction there is an associated Participant agent on the server. The Participant agents associated to a given user directly report to, and eventually get orders from, his or her Personal Agent.

Initiator Participant is a special Participant agent that represents the user, with the role to create and initiate the auction. For example, in an English auction the Initiator Participant represents the user that has the role of a seller, while the remaining Participants represent users that have the role of buyers. Usually, when a user initiates an auction he or she can specify also a condition that triggers the start of the auction (if this condition is missing then it is assumed to be true by default, which means that the auction will start immediately after creation). Examples of starting conditions are: predefined starting time or registration of a minimum number of participants (this second condition is particularly useful for setting up experimental scenarios, see Section 5).

### 3.1.2 Core Layer

The core of the server is represented by Auction Service, Auction Manager, Auction Directory, and Proxy agents. This layer is responsible for the auctions’ management and for the coordination of the auction participants by implementing the rules that govern the auction.
The *Auction Service* is the agent that manages all the active auctions registered with the server. This agent is the entry point of the auction service and it is responsible for creation of new auctions, as well as with registering of new participants to an active auction.

*Auction Manager* manages a single active auction on the server (also known as auction instance in [2]). The *Auction Manager* coordinates the participants registered to that active auction. There is a separate *Auction Manager* agent for each active auction in the system. It implements a specific type of auction – English auction in this case – but in principle an *Auction Manager* can be configured to support an arbitrary auction type. The management of active auctions includes the activities that usually occur in an auction, i.e.: auction creation, bidding, agreement formation and auction termination. The *Auction Manager* is created by the *Auction Service* when a user wants to sell a product through an English auction. The *Auction Manager* has the following responsibilities:

- to request the creation and destruction of *Proxy* agents
- to request the creation of *Participant* agents and to assign them to the right *Proxy* agent
- to notify all the *Proxy* agents when the value of the highest bid was updated
- to accept or reject the bids from *Participant* agents that are forwarded by *Proxy* agents
- to manage the parameters of the auction, including: auction name, starting price, product name, starting and ending dates, as well as the status of the auction (for example, the currently highest bid)
- to trigger the auction termination when the time has expired by informing all the *Proxy* agents that the auction has finished
- to record the winner and the final price of the auction.

*Auction Directory* agent manages the registry of active auctions, as well as the identifiers of their associated *Auction Manager* agents. Potential auction participants can search through this registry to find an active auction that meets their requirements.

Each *Proxy* agent handles the bids received from a subset of *Participant* agents. The *Participant* agents are split into disjoint groups, and each group is managed by a single *Proxy* agent. *Proxy* and *Participant* agents are linked into a balanced two-level hierarchical structure rooted at the *Auction Manager* such that the total number of *Proxy* agents is at most equal to the number of *Participant* agents that are linked to each *Proxy* agent. Using this balancing criterion we can reduce the time for processing incoming bids. During bidding, *Participant* agents communicate heavily with their *Proxy* agents, while *Proxy* agents pass on to the *Auction Manager* agent only the relevant bids and the other bids are filtered out and processed locally, thus reducing the amount of messages handled by the *Auction Manager* that results in enhancing the server response time. Note that this is particularly important in
the assumed model of an English auction (deadline-driven), where a large number of bids can materialize near the auction-deadline (due to the snipping).

Proxy agents form an intermediate layer between the Auction Manager and the Participant agents. Their main responsibility is to take over a part of the load that is necessary for bid processing that was initially the sole responsibility of the Auction Manager [2]. Each Proxy agent records a local currently highest bid (that can be different from the currently highest bid of the auction recorded by the Auction Manager) and updates it regularly based on the notifications received from the Auction Manager and bids received from Participant agents. The responsibilities of the Proxy agent are:

• to filter out the bids received from the Participant agents passing up to the Auction Manager only those bids that meet the improvement rule, i.e. are higher than the local currently highest bid known by the Proxy
• to notify Participant agents about the acceptance or rejection of their bid
• to receive notifications about the update of the currently highest bid from the Auction Manager
• to notify Participant agents after receiving such an update from the Auction Manager; the local currently highest bid of the Proxy is also updated.

Let us suppose that, on average, a Proxy agent is managing \( k \) Participant agents and that the total number of Proxy agents is \( p \). This means that the total number of participants is \( n = p \times k \). According to our coordination model of an English auction, whenever a new bid is accepted by the Auction Manager, all the participants must be notified accordingly, and this process will obviously require a time \( O(n) \). However, with the new Proxy-based hierarchical approach the time will be \( O(p+k) \).

This value can be optimized to \( O(\sqrt{n}) \) for \( n \) Participant agents if we set the constraint that the number of participants managed by each Proxy agent will never outnumber the total number of Proxy agents connected to the Auction Manager, while the number of Proxy agents is increased whenever this is really necessary to manage all the registered Participant agents. Basically this means that \( p \sim k \) from which we can derive the \( O(\sqrt{n}) \) average time needed to process a bid.

Proxy and Participant agents are linked into a balanced two-level hierarchical structure rooted at the Auction Manager such that the maximum number of Participant agents that are linked to each Proxy agent is at most equal to the total number of Proxy agents. The dynamic creation and destruction of Proxy agents will take into account the preservation of this balancing requirement.

An example showing the details of the relationship between the Auction Manager, the Proxy agents, and the Participant agents is presented in Figure 2. This figure shows 5 Participant agents \( p_1, \ldots, p_5 \) registered to the same auction. Let us assume that the participants registered to the auction in this order. When participant \( p_5 \) registered, only 2 Proxy agents \( x_1 \) and \( x_2 \) were present on the server. According to the balancing criterion, the registration of the new Participant agent \( p_5 \) triggered the creation of a new Proxy agent \( x_3 \), as well as the link of \( p_5 \) to \( x_3 \).
The new $p_5$ Participant agent cannot be assigned to one of the existing Proxy agents because that would mean that at least one of them would have 3 Participant agents, while the number of Proxy agents is 2. So a new Proxy agent $x_3$ must be created.

Requests for creating and destroying Proxies are issued by the Auction Manager. A Proxy agent is created when a new Participant is created and $n > p^2$ where $n$ is the total number of Participants including the newly created one, and $p$ is the number of Proxies.

When a Participant quits an auction, if $n \leq (p - 1)^2$ then the Proxy with the smallest number of participants is deleted, where $n$ is the total number of participants (after the removal of the current Participant from the server) and $p$ is the current number of Proxies. The Participant agents that were linked to the Proxy that will be deleted will be evenly reallocated to the remaining Proxy agents with the smallest number of Participants.

Finally, there is an issue concerning what happens with the Proxy agents when the auction terminates. An English auction is running for a certain time duration, which is set when the auction is created. Optionally, this duration can be extended by the Auction Manager with a short period – timeout, in order to avoid problems created by “late bidding”, by allowing all interested bidders to submit their bids. This is a solution used by auction servers to increase their “fairness” [5].

The timing of the auction is controlled by the Auction Manager agent. When this agent detects and declares the auction terminated, it will inform the Proxy agents about the auction termination and the auction result. Consequently, Proxy agents will reject forthcoming bids submitted by non-acceptably late bidders. Moreover, Proxy agents will notify their Participant agents accordingly, while Participant agents notify in turn their Personal Agent. Finally, Participant agents self-destroy. This process can optionally trigger up in the tree the self-destroy of Proxy agents, to maintain the tree balanced.
Note that the process for managing the creation and destruction of Participant and Proxy agents can be improved using the technique of managing resource pools – in this particular case we would have agents pools on each computer of the network [40].

3.1.3 Resource Layer

The agents Computer Manager and Resource Manager constitute the part of the server that is responsible for the management of the computational resources, which run the server software.

When the server is installed on a computer network, the Computer Manager agents are responsible for the basic management (i.e. creation/destruction) of agents on each available machine. There is one Computer Manager agent for each computer that is part of the auction server infrastructure. This agent keeps track of all the Participant and Proxy agents that were created and are active on that computer. A Computer Manager agent gets requests from the Resource Manager agent that contain the agent name and type and then creates the agent on that computer accordingly.

Resource Manager contains a registry of all the agents of type Computer Manager from the system. More exactly, when the Auction Manager decides to create a new Participant or Proxy agent, it will ask the Resource Manager to perform this operation. Currently, the Resource Manager agent keeps track of the number of agents already created on each available machine. On the basis of this information, it decides where should the new required agent be created and orders the creation accordingly to the corresponding Computer Manager agent. However, in a more general setting we can expand the functionality of the Computer Manager agents to monitor the resources and the network load of each single computer of the server. Then, the Resource Manager would be able to query Computer Manager agents before deciding where (on what computer) to order the creation of new agents. This extension was left as future work.

3.2 Agent Interactions

Let us now summarize the protocols for interacting with the core of the auction server. Basically, they follow the same rules as the initial proposal of [3] that we also considered in [2]. According to these protocols, a user represented by his or her Personal Agent can create auctions, subscribe to active auctions, submit bids, receive replies about bid acceptance or rejection, receive notifications about the update of an auction status, and receive notifications about auction termination and auction winner.

Note that the interaction protocols presented and discussed in this section are focused on what is happening “inside” the auction server. However, the management of the link between a user and his or her Personal Agent is realized “outside” the auction server – see Section 4.
A new auction is created when a Personal Agent sends to the Auction Service a createAuction request (see Figure 3). Then the Auction Service creates an Auction Manager that represents the new auction. The Auction Manager orders creation of a new Initiator Participant agent linked to the Personal Agent of the user seller, as well as of a new Proxy for this participant, to the Resource Manager. The Auction Service then confirms to the Personal Agent the creation of the Initiator Participant agent. The auction description is also added to the Auction Directory. The Initiator Participant and Auction Manager agents will further interact during the auction process.

A Personal Agent can query the Auction Directory about active auctions (see Figure 4). The Personal Agent then chooses an auction and contacts the Auction Manager of that auction which should create a Participant for it.

The Auction Manager decides to create a new Participant agent and, optionally, a new Proxy agent. Their creation is handled by the Resource Manager agent, while the actual creation is performed by the Computer Management agent. The Auction Manager requests the Resource Manager agent to choose an appropriate computer where these agents will be created. The Resource Manager determines this computer and instructs the corresponding Computer Manager agent, to perform the creation action. When a computer reached the maximum acceptable load another computer should be used. If all computers are fully loaded then the computer with the smallest number of proxies or participants (in that order) is chosen; however, it is expected that the performance of the server will degrade in such situations. All Participant agents are created on the same computer as their Proxy agent.

The message exchanges that are needed for creating Participant and Proxy agents are shown in Figure 5.
Each Participant agent sends his or her bids to their designated Proxy according to the protocol shown in Figure 6. This bid is then compared by the Proxy to the currently best bid. Higher bids are reported to the Auction Manager while lower bids are rejected with a refuseBid message. Then the Auction Manager saves the highest bid (bestBid state variable) and refuses submitted bids that are lower than the bestBid. The refusal message is then returned by the Proxy agent back to each originating Participant that did not submit the current bestBid. If a bid higher than bestBid is received then the Auction Manager responds to the Proxy with a bidAccepted message, which is propagated to the originating Participant. If the bestBid value has changed then the Auction Manager informs all the Proxy agents. Then each Proxy will propagate the value of the new bestBid to its Participant.
agents. This forwarding is represented in Figure 6 as a multicast message that is displayed as a little circle at the end of the message arrow.

Figure 6. The bid submission protocol

4 INTEGRATION INTO AN ONLINE AUCTION SYSTEM

In this section we outline the design of an online auction system that incorporates the Agent-based Auction Server. The system design is composed of three parts: system architecture, interaction protocols, and design details of the system components.

4.1 System Architecture

The auction server was implemented using the JADE agent platform [10]. For experimenting with the usability of this auction server we have developed an online system equipped with a Web-based GUI that allows human users to create, search for and participate in English auctions. Therefore, the architecture will contain a special subsystem dedicated to the interface of JADE agent middleware with the Web-server.
The architecture of the application follows the classical separation between the client side, comprising the human user equipped with a Web browser, and the server side comprising a Web server that interacts with the auction server. Therefore we designed and developed a software for binding the Web server part that is non-agent software with the auction server that consists solely of JADE-based agent software.

![System architecture diagram](image)

**Figure 7. System architecture**

The system has a multi-layer architecture composed of the following layers, as shown in Figure 7:

**User layer.** This layer represents the client part of the system that consists of a Web browser combined with HTML content including JavaScript code that is downloaded from the Web server.

**Web layer.** This layer supports the user interaction functionalities. It consists of a Web server enhanced with a set of Java servlets that implement the user functionalities of the online application.

**Binding layer.** This layer is represented by the software that enables the interfacing of the non-agent Web server software with the Agent-based Auction Server. This software is encapsulated into a special servlet called Agent servlet that is able to communicate with the JADE platform.
Agent layer. This layer represents the Agent-based Auction Server that was built on top of JADE platform.

4.2 Design Details

4.2.1 User and Web Layers

The Web layer is responsible for management of users and their accounts, while the auction server (i.e. the Agent layer) is responsible for auction management. With this separation of functionalities, the Web layer will support the interaction of the application with the human user, via a Web-based GUI that is based on HTML, Asynchronous JavaScript, and XML (i.e. Ajax [35]).

As we did not create a functionality for user authentication in the Agent-based Auction Server, we had to provide a solution for this problem at the level of the Web layer. So, in our prototype system, the Web layer is responsible for user authentication (login and logout functionalities) and management of user accounts. The addition of this functionality requires the Web layer to maintain a separate database for the management of user account information.

At the User layer, the servlets must provide HTML responses with information extracted from the Agent layer using software from the Binding layer. This operation can be time consuming and thus it can slow down the load time of the Web page. This issue was addressed by inserting JavaScript code into the HTML responses. This code allows the Web page to update itself quickly and efficiently using asynchronous requests. The JavaScript code issues automatic or user generated HTTP requests to the Agent servlet which was configured to reply with the XML responses. Note that the Agent servlet is the only servlet of the Web layer that can communicate with the agents on the Agent layer (via the Binding layer).

4.2.2 Binding Layer

The Binding layer communicates with the Agent layer using a specialized software. This software benefits from the JADE facilities for interfacing agent and non-agent software materialized as the JadeGateway class ([22]). The interface is achieved using agents (also known as Gateway agents) that are created locally by the Agent servlet. One Gateway agent is created for each user logged into the system. These local agents are created in a local container that is connected to the agent platform that hosts the auction server. This container is created and started together with the Web layer.

Whenever a new user logs into the system, the Agent servlet automatically creates a local Gateway agent with the role of relaying messages from the Web layer to the agents on the auction server. Then the Agent servlet locally creates and passes a serializable object (called “blackboard object” in [22]) to the Gateway agent assigned to the current user. The Gateway agent then sends the message using the JADE messaging functionality to an agent located on the auction server,
Figure 8. User login

according to one of the interaction protocols presented in Section 4.3. This type of interaction is marked with the $\langle \text{object2agent} \rangle$ stereotype in Figures 8, 9, 10, 11, and 12. Conversely, whenever an agent of the auction server must send information to the Web layer it will use the JADE messaging to send this information to the corresponding Gateway agent located in the Agent servlet. Then the Gateway agent will invoke a method to update the “blackboard object” thus achieving the correct transfer of the information to the Agent servlet. This type of interaction is marked with the $\langle \text{agent2object} \rangle$ stereotype in Figures 10 and 11.

Figure 9. Initiation of an auction

Note that whenever the JavaScript code of the Web page asynchronously requests an information update from the Agent servlet the servlet will respond with a message containing the relevant data extracted from the “blackboard” object and
represented in the XML format. This type of interaction is represented with the ≪XML response≫ stereotype in Figures 10 and 11.

![Diagram of searching and joining an auction](image)

**Figure 10. Searching and joining an auction**

### 4.3 Interaction Protocols

In this section we formally describe the interactions that happen between the software components of our auction system, “outside” the auction server, as opposed to the agent interaction protocols described in Section 3 that are focused on the “inside” of the auction server. Please note that although some of these interactions are related to the same activity – for example the bidding activity has a part inside the server, as well as a part outside the server, they are presented separately (in Section 3 and Section 4) for at least two reasons:

1. their common part is minimal (it is reduced to the *Personal Agent*) and they can be well-understood separately;
2. the auction server is a separate subsystem that can be integrated in other types of applications, for example using a Web-service interface.

The *user login* operation is detailed in Figure 8. The first part (interactions numbered from 1 to 7) achieves the authentication function. If the authentication is successful (i.e. the interactions proceed according to the branch consisting of messages 6 and 7) the *Agent-based Auction Server* is notified accordingly, via the interaction consisting of messages 8 and 9. Message 8 is a notification sent to the *AgentServlet* that the user logs into the system. Consequently, the link between the
user name and the identifier of his or her *Personal Agent* is retrieved. Eventually, the *Personal Agent* of the user must be created and started (message 9), either if the user logs into the system for the first time or if his or her *Personal Agent* was offline.

The operation of *initiating an auction* is detailed in Figure 9. The first 3 interactions activate the user menu for setting the auction data. The next 3 interactions support the function of creating a new auction. The actual creation is achieved after interaction 6. Note that we assume that when a new auction is created the user is already logged in and its *Personal Agent* is active.

The operation of *searching and registering at an active auction* is detailed in Figure 10. The first 3 interactions activate the user menu for setting the search criteria for the desired auction. The next 6 interactions (numbered from 4 to 9) support the function of searching auctions available in the auction directory. This is achieved with the help of the *Auction Directory* agent, residing on the auction server. Then the user chooses the desired auction (this is achieved by interaction 10). Note that after this action the name of the corresponding *Auction Manager* agent becomes known. Finally, the last 3 interactions (numbered from 11 to 13) allow the user to join the desired auction.

The operation of *bidding in an auction* is detailed in Figure 11. The first 3 interactions activate the user menu for setting the bid data. The next 3 interactions (numbered from 4 to 6) support the function of submitting the bid to the auction server. Note that we assume that at this point the user knows the identifier of the
auction where he or she wishes to submit the bid (parameter auction). Interaction 7 is happening whenever the Personal Agent is notified by the Agent Layer that at least one of the auctions where the user is registered has updated its status. The last 2 interactions (8 and 9) allow the user to visualize the quotes of the auctions where he or she is subscribed. In particular, for an English auction, the user can check if a given bid was accepted or not, by the auction server. The updates are periodically triggered by a timer incorporated into the JavaScript code that runs in the user’s browser.

The userlogout operation is detailed in Figure 12. Similarly to the login operation, the auction server is notified that the user is leaving the system. However, note that in this case the notification is sent directly to the Personal Agent that represents the user.

Finally, it is important to observe that there is an interesting relationship (not shown in Figures 8 and 12) between the Auction Service and the Personal Agent that represents a specific human user on the auction server. This is the result of the fact that the Personal Agent has a very important role, by controlling the user participation in auctions, even when the user is disconnected from the online system. This fact has two important consequences:

- During the login operation the system must check if the user already has an active Personal Agent at the auction server, and if not it must create one. We assume that this operation is achieved by the Auction Service agent. So, the Auction Service agent has the responsibility of creating and setting up of a new Personal Agent according to user requirements. An example of such requirements could be: what to buy, the maximum price and the acceptable auction duration. Note that such requirements will become goals of the Personal Agent.

- During the logout operation the system must inform the Personal Agent that the user has left the system. However, the Personal Agent can behave more or less autonomously (according to the user requirements) in representing the user preferences and interests. So, the Personal Agent can autonomously decide to go offline in situations when, for example, there are no more active auctions in which the user is participating or, more generally, when a certain user goal was either achieved or is considered not achievable given the current state-of-affairs. Alternatively the Personal Agent can decide to continue its execution on the server.

Nevertheless, we set the requirement that the Personal Agent must always notify the Auction Service that it will go offline before doing so, such that if the user logs onto the system again then the Auction Service will be able to create and setup a new Personal Agent accordingly. The Personal Agent will autonomously decide to go offline whenever there is nothing left to do for the user. In particular, this could happen when the user logouts and he/she is not currently involved in any auctions, as well as if there are no active goals of the Personal Agent agenda to be pursued. This can happen for example either when all the auctions where the user
was involved are finished or when the user just logged in for the first time, did not create any auctions and did not set any requirements for the Personal Agent but it just decided to leave the system, i.e. to log out.

Figure 12. User logout

5 EXPERIMENTS AND USABILITY

5.1 Experiments with the Distributed Agent-Based Auction Service

We experimentally evaluated the current architecture by comparing it with our initially proposal from [2], where no Proxies were used. Then we conducted initial scalability experiments by running our system on 2 and 3 computers.

For the experiments we used a network of dual core processors at 2.5 GHz and 1 GB of RAM memory. These workstations were interconnected using a high-speed Myrinet interconnection network at 2 Gb/s. According to [37], “Myrinet is a [...] high-performance, packet-communication and switching technology that is widely used to interconnect clusters of workstations”. As multi-agent middleware platform we have used JADE 4.0 [38].

In this experiment the participants are allowed to bid automatically, so they were equipped with a bidding strategy to tell them if, when, and how (much) to bid. The bidding follows a snipping scenario, i.e. as soon as a participant receives a notification that s/he was outbid by another participant, s/he immediately submits a higher bid by adding a predefined increment to the value of the currently highest bid. During the snipping scenario the auction server is heavily loaded with bid processing activities.

The starting price of the auction was set to 0. The auction duration was set to 1.5 minutes in all the experiments. The increment value for the participants’ strategy was set to 10. The agents were allowed to bid up to a very high reserve price (100 000). The auction terminates when its allocated time expires. In our cases, to assure that we actually study the performance of the server, as the reserve prices were set to very high values, the auctions end before any of the agents reaches their reserve price.
Note that whenever a Participant Agent bids according to this strategy, it must receive an answer confirming if the bid was successful (or that it was rejected). This is very important, as we are in a distributed environment with multiple agents bidding concurrently, and it might happen that even if an agent is choosing a high enough value to bid, it might be outbid by another agent that submits its bid almost simultaneously. The following performance measures were recorded in our experiments:

- **Latency** = the average time it takes the system to answer a bid.
- **Throughput** = the number of bids handled per unit of time; this value is calculated by dividing the total number of answered bids by the duration of the auction.

The setup of the experiment assumes running a script that starts the JADE multi-agent platform and automatically creates the *Auction Service* and *Auction Directory* agents. Then the *Personal Agents* are created for each user that participates in the auction.

The *Personal Agent* that initiates the auction was configured to set the condition for starting the bidding when a specific number of participants has joined the auction. This condition is configured into the *Auction Manager* that governs the auction.

When a certain given number of participants is reached, this *Auction Manager* will enable the starting of the bidding process. Basically, with this approach we were only looking for a simple method to set up our experiment consisting of many agents bidding aggressively in an auction, while keeping the consistency with the design philosophy of the auction server.

We ran our experiments using an increasing number of *Personal Agents* and we calculated the performance measures by running the framework on one, two and three computers. In order to also compare with our prior approach we also ran a version of the program without proxies, forcing the *Auction Manager* to handle all bids. In each test we ran only one auction. The case when multiple auctions are run in parallel was left as future work.

The results of our experiments are shown in Figures 13 and 14, as well as in Tables 1 and 2.

<table>
<thead>
<tr>
<th>Agents</th>
<th>No Proxy</th>
<th>1 Comp.</th>
<th>2 Comp.</th>
<th>3 Comp.</th>
</tr>
</thead>
<tbody>
<tr>
<td>500</td>
<td>5087.53</td>
<td>2172.14</td>
<td>148.55</td>
<td>104.39</td>
</tr>
<tr>
<td>1000</td>
<td>20038.9</td>
<td>8114.66</td>
<td>266.85</td>
<td>217.01</td>
</tr>
<tr>
<td>1500</td>
<td>39752.5</td>
<td>8781.39</td>
<td>426.50</td>
<td>277.88</td>
</tr>
<tr>
<td>2000</td>
<td>66852.3</td>
<td>11691.2</td>
<td>575.06</td>
<td>452.73</td>
</tr>
</tbody>
</table>

Table 1. Latency [ms]

Note that when we ran our system with 2000 bidding agents without proxies only 1273 agents actually got to bid at least once during the allocated time of 1.5 minutes.
5.2 Usability of the Online Auction System

We now consider a sample use case involving four users $U_1$, $U_2$, $U_3$, $U_4$ participating in auctions with the help of our system. On this use case we highlight the usability aspects of the system, as well as some of the details regarding the information exchanged by the various components. In particular, we are interested in checking the information flow that is triggered into the system by user initiated actions. Please note that in this description we will make references to the diagrams introduced in Section 4.

<table>
<thead>
<tr>
<th>Agents</th>
<th>No Proxy</th>
<th>1 Comp.</th>
<th>2 Comp.</th>
<th>3 Comp.</th>
</tr>
</thead>
<tbody>
<tr>
<td>500</td>
<td>0.141</td>
<td>1.115</td>
<td>1.527</td>
<td>1.801</td>
</tr>
<tr>
<td>1000</td>
<td>0.065</td>
<td>0.949</td>
<td>1.714</td>
<td>1.744</td>
</tr>
<tr>
<td>1500</td>
<td>0.035</td>
<td>0.851</td>
<td>1.43</td>
<td>1.86</td>
</tr>
<tr>
<td>2000</td>
<td>0.025</td>
<td>0.779</td>
<td>1.502</td>
<td>1.672</td>
</tr>
</tbody>
</table>

Table 2. Throughput [no. of bids/ms]
The users log into the system at the *User Layer* using the *Login GUI* in the Web browser shown in Figure 15. This operation, shown as message 1 in Figure 8, is needed to communicate the username and password to the *Web Layer*. The login request reaches the *Login Servlet* (message 2 in Figure 8 – an *Http Request* message). The user credentials are then verified and an HTML page is returned to the user presenting one of the following two possible outcomes:

- If the credentials are incorrect then an HTML page reports an error to the user’s browser (message 5 in Figure 8).
- Otherwise the HTML page contains JavaScript code for interacting with the *AgentServlet* (message 7 in Figure 8). After the interaction with the servlet (message 8 in Figure 8), a *Personal Agent* is created for each user.

In our sample use case, four *PersonalAgents* (*PA*₁, *PA*₂, *PA*₃, *PA*₄) will be created in the *Agent Layer*.

We assume that user *U*₁ creates two auctions labeled *A*₁ and *A*₂ (their details are shown in Table 3). For each auction, participant (initiators in this case) *IPart₁, IPart₂* agents, as well as *AuctionManager₁* and *AuctionManager₂* agents are created.

Interaction 1 in Figure 9 represents the initial request sent by *U*₁ for creating an auction. Interaction 2 in Figure 9 represents the *HTTP Request* message sent to the *AddAuction Servlet*. This servlet responds with an HTML file containing
also JavaScript code for creating the form Create Auction GUI shown in Figure 16. The form enables the user to input the product descriptions as shown in Table 3 (message 4 in Figure 9). These are then passed to the AgentServlet using message 5 in Figure 9. Finally, AgentServlet requests PA<sub>1</sub> agent (using message 6 in Figure 9) to create the auction in the Agent Layer.

<table>
<thead>
<tr>
<th>auction</th>
<th>A1</th>
<th>A2</th>
</tr>
</thead>
<tbody>
<tr>
<td>product name</td>
<td>Desktop</td>
<td>TV</td>
</tr>
<tr>
<td>description</td>
<td>I3, 4 GB RAM, 500 GB HDD</td>
<td>LCD, 81 cm</td>
</tr>
<tr>
<td>starting price</td>
<td>1500 RON</td>
<td>800 EUR</td>
</tr>
<tr>
<td>increment</td>
<td>10 RON</td>
<td>5 RON</td>
</tr>
<tr>
<td>auction duration</td>
<td>3 h</td>
<td>5 h</td>
</tr>
</tbody>
</table>

Table 3. Auction details

We now assume that users U<sub>2</sub>, U<sub>3</sub> and U<sub>4</sub> are searching for a desktop computer which is provided by auction A<sub>1</sub>, while user U<sub>3</sub> is also looking for a TV which is provided by auction A<sub>2</sub>. Searching for a GUI starts when a user issues a request that triggers messages 1 and 2 in Figure 10. The SearchServlet responds with an HTML form that represents the Join Auction GUI (see Figure 17). Then the user communicates his or her search parameters to the system (message 4 in Figure 10). The search parameters are passed to the Auction Directory using messages 5, 6 in Figure 10. The Auction Directory replies by triggering the sequence of messages 7, 8 and 9. The result contains the description of a list of auctions as an XML file (more precisely message 8 in Figure 10). The user can choose the auctions they want to join by triggering interactions 10, 11 and 12. Whenever a user joins an auction, a new Participant agent is created to represent the user bidding in that auction (this is achieved via message 13 in Figure 10). In our scenario Part<sub>2</sub>, Part<sub>3</sub>, and Part<sub>4</sub> agents are created to represent users U<sub>2</sub>, U<sub>3</sub> and U<sub>4</sub> acting in auction A<sub>1</sub>, as well as Part<sub>5</sub> agent is created to represent user U<sub>3</sub> acting in auction A<sub>2</sub>.

![Figure 17. Search and Join Auction GUI](image-url)
In what follows we shall focus only on what happens in auction $A_1$. Let us assume that user $U_2$ bids 1510 RON and his/her bid is accepted. Message 5 in Figure 11 contains the bid information that user $U_2$ sends to the AgentServlet: the bid value 1510 RON and the participant ID Part2. The bid value is input by the user via the GUI shown in Figure 18. The AgentServlet determines the identifier of the PersonalAgent attached to user $U_2$, i.e. $PA_2$ and then forwards the bid information to $PA_2$ (message 6 in Figure 11). Whenever a PersonalAgent receives a notification about the update of the currently highest bid, it notifies the AgentServlet (message 7 in Figure 11). The GUIs of all the users participating in auction $A_1$ (i.e. $U_2$, as well as $U_3$ and $U_4$ via their $PA_3$ and $PA_4$ agents) are automatically updated about the currently highest bid by the JavaScript code that periodically retrieves from the AgentServlet the updated information encoded an XML message (message 9 in Figure 11). In this example the XML message contains the information shown in Table 4.

<table>
<thead>
<tr>
<th>ProductName</th>
<th>Description</th>
<th>LastBid</th>
<th>Winner</th>
<th>Bid</th>
<th>TimeLeft</th>
</tr>
</thead>
<tbody>
<tr>
<td>Desktop</td>
<td>i3 3.50M Processor, 4 GB RAM, 500 GB HDD</td>
<td>1510.0</td>
<td>Part2</td>
<td>1510</td>
<td>0 D 0:27:43</td>
</tr>
</tbody>
</table>

Table 4. Currently highest bid update

In what follows let us assume that user $U_2$ logs out. However, his or her PersonalAgent, i.e. $PA_2$ is kept alive on the Agent-based Auction Server and continuously receives updates as auction $A_1$ is proceeding. Now, assuming that $U_3$ and $U_4$ both bid 1520 RON, with $U_3$ being slightly faster than $U_4$, the bid of $U_3$ is accepted, while the bid of $U_4$ is refused. Now, as $U_4$ is not happy that his or her bid was rejected, $U_4$ will submit a new higher bid of 1530 RON. At this point $U_3$ resigns the auction.

Let us now assume that $U_2$ logs in again. It will be automatically informed by $PA_2$ agent that the currently highest bid is 1530 RON and it was submitted by $U_4$, as shown in Figure 19.

Now, let us assume that $U_2$ decides to place a new bid of 1540 RON, but meanwhile the auction time expired. The bid submitted by $U_2$ will be ignored by the
server. The server notifies all the users, including $U_2$, about the outcome of auction $A_1$. Assuming that auction $A_2$ is also finalized without any winner (more exactly, no bids were submitted in this auction), $U_2$ will receive the notification shown in Figure 20.

![Figure 19. New updated highest bid](image1)

![Figure 20. Final notification](image2)

6 CONCLUSIONS

In this paper we described in detail an improved agent-based architecture for an English auction server. The initial experimental results show that our hierarchical scheme of structuring the server using Proxy agents and a simple balancing scheme is effective and has good scalability, when the server is distributed on multiple machines. As future work we plan to:

1. strengthen the results by performing more experiments on larger networks; in particular we will target experiments on high-performance computer clusters;
2. extend the architecture to other types of auctions.

This paper also introduced our design and implementation of an online auction system that incorporates the Agent-based Auction Server. The system provides a Web-based GUI for the Agent-based Auction Server. We outlined the main functionalities of the system, as well as their design and implementation, in terms of system architecture, design details and interaction protocols. The main outcome of our work is a clean specification of the Web-based and agent-based software layers of our system, as well as of their software interfaces. As future work we plan to:

1. expand our design by providing a Web services interface to our Agent-based Auction Server;
2. investigate the relation between the human user and his or her Personal Agent, in particular on how human requirements in the area of auctioning and e-commerce can be mapped onto elements of the Personal Agent architecture, like for example those related to the BDI model.

REFERENCES


Costin Bădică is Professor at the Department of Computers and Information Technology, University of Craiova from 2006. In 2001 and 2002 he was Postdoctoral Researcher at the Department of Computer Science, King’s College, London, United Kingdom. His research interests are at the intersection of artificial intelligence, distributed systems and software engineering. He was involved in many national and international research projects. He is coordinating the “Intelligent Distributed Systems” research group (http://ids.software.ucv.ro). He co-authored more than 100 research publications, coedited books, book chapters, editorials of journal special issues. He is member of the editorial board of several international journals and served as member of the program or organizing committee of many international conferences.

Sorin Ilie is research assistant at the Department of Computers and Information Technology, University of Craiova from 2009. He defended his Ph.D. thesis in December 2012 on Distributed Ant Colony Optimization. His research interests include artificial intelligence, distributed systems and software engineering. He was involved in four national and international research projects. He is a member of the “Intelligent Distributed Systems” research group (http://ids.software.ucv.ro). He co-authored 28 research publications. He served as member of the program or organizing committee of two international conferences.

Alex Muscar is a Ph.D. student at Department of Computers and Information Technology, University of Craiova from 2010. Before starting his Ph.D. studies, he developed large scale applications in Python on Google’s AppEngine cloud platform and high performance recommender systems in C++. His research interests are mainly in the areas of programming languages and concurrent and distributed systems. He is also a member of the “Intelligent Distributed Systems” research group.
Amelia Bădică holds a Ph. D. in Economics and she currently works as Senior Lecturer at the Business Information Systems Department, University of Craiova, Romania. She has a specialization in management information systems at Binghamton University, USA. Her research interests cover the application of expert systems, software engineering and Web technologies in business and management. She authored many papers in journals and conference proceedings on these subjects. She was involved as principal investigator in a research project concerning data extraction from the Web. She also served as member of the program or organizing committee of many international conferences.

Liviu Sandu has a Bachelor degree in Computer Science from the Faculty of Automatics, Computers and Electronics, University of Craiova, Romania. Now he is studying for the Master’s degree at Faculty of Automatics, Computers and Electronics, University of Craiova, Romania. From 2011 he also works as a software developer at RT-SOF T SRL, Romania.

Raluc a S b o ra has a Bachelor degree in Computer Science from the Faculty of Automatics, Computers and Electronics, University of Craiova, Romania. Now she is studying for the Master’s degree at Faculty of Automatics, Computers and Electronics, University of Craiova, Romania. From 2012 she also works as a Software Tester at Hella, Romania.

Maria G an z h a is Assistant Professor, obtained her M. Sc. and Ph. D. in Applied Mathematics from the Moscow State University, Moscow, Russia in 1987 and 1991, respectively. Her initial research interests were in the area of differential equations, solving mixed wave equations in space with disappearing obstacles in particular. Currently she works in software engineering, HPC, distributed computing and agent systems in particular. She has published more than 130 research papers and is on editorial boards of 10 journals and a book series and was invited to Program Committees of over 150 conferences.
Marcin Paprzycki has received his M.Sc. Degree in 1986 from Adam Mickiewicz University in Poznan, Poland, his Ph.D. in 1990 from Southern Methodist University in Dallas, Texas and his Doctor of Science Degree from Bulgarian Academy of Sciences in 2008. His initial research interests were in high performance computing and parallel computing, high performance linear algebra in particular. Over time they evolved and started to include also distributed systems and Internet-based computing; in particular, agent systems. He has published more than 300 research papers and was invited to Program Committees of over 400 international conferences. He is on editorial boards of 16 journals.